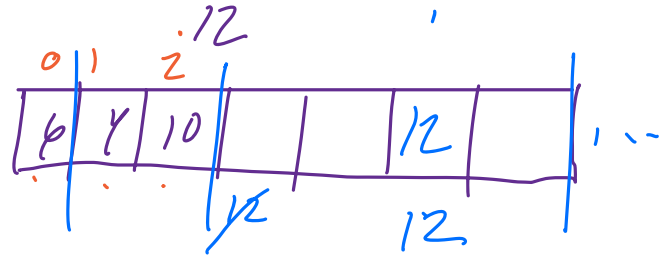
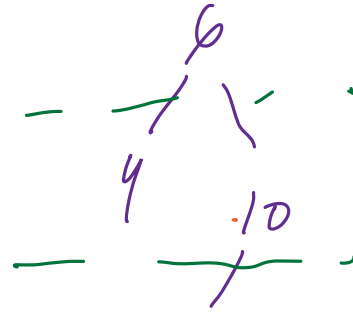
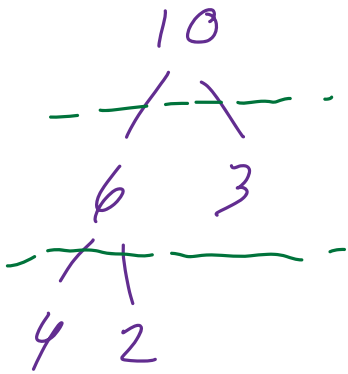


Recap: Embedding trees/heaps in Arrays

WHAT ARE THE ARRAYS FOR THESE TREES?



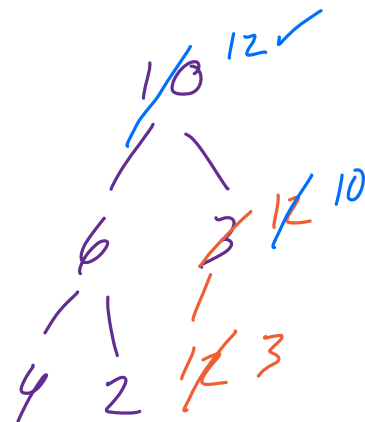
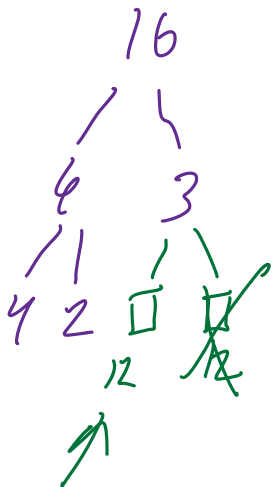
- For some node at index i :
- $\text{left}(i)$ is at $(i * 2) + 1$
 - $\text{right}(i)$ is at $(i * 2) + 2$
 - $\text{parent}(i)$ is at $\text{floor}((i - 1)/2)$

Why not keep stuff condensed? Need to respect parent/child formulas to match the structure of the tree
 => Without spaces, parent, left/right formulas will break

How to insert?

In the code version: always add new elements to the next available space
 => Keeps balanced
 => Always know how to find next slot (just the end of the array)

Ex. INSERT 12

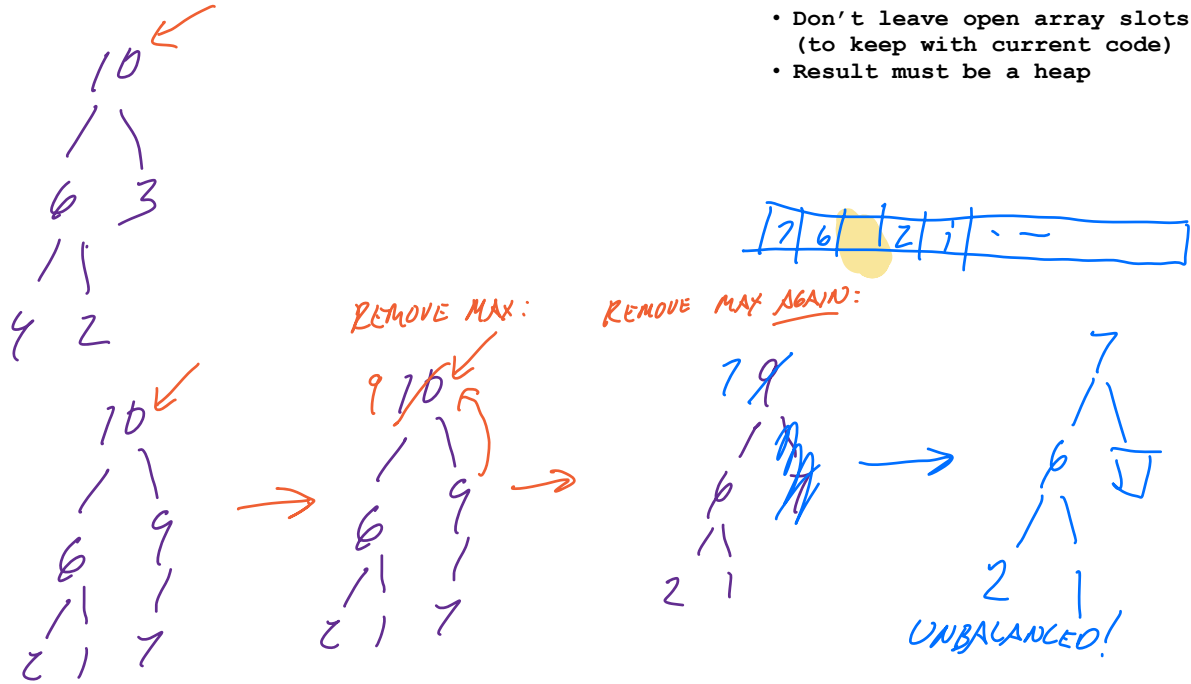


How would we implement `delete_max`?

Let's try something similar to `insert`, where we swap nodes up after removing an element:

Requirements

- Maintain balance
- Don't leave open array slots (to keep with current code)
- Result must be a heap

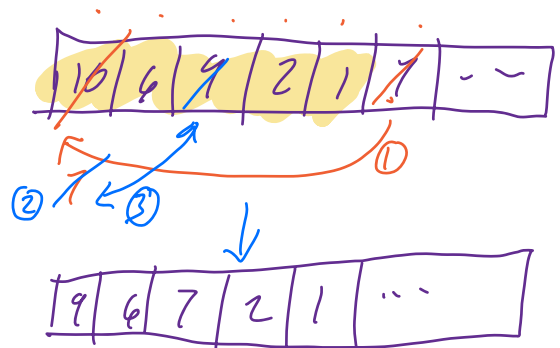


Problem: If use the same remove + swap from here, can end up with unbalanced tree
 => Need to do delete without creating gaps
 => If we can just swap values within existing cells in use, can avoid gaps
 => Leverages array structure!

We didn't implement `remove_max`, but here's the intuition:

- If we're removing an element, the resulting array must be one smaller than before
- Removing the max element creates a hole => swap the last element (ie, the one that would get "abandoned" if we were to shrink the array) to the top (in the hole left by the max)
- Swap this new top element down until the result is a heap

This idea leverages the same principles as `insert`, but involves swapping in a different direction. Similar to how we knew where to add a new element when inserting into the heap, we leverage the array to know where to find a new element to start swapping down (ie, last one)



- ① REMOVE 10
- ② SWAP LAST (7) TO FRONT
- ③ SWAP 7 DOWN UNTIL RESULT IS HEAP

RECAP: NOW TO THINK ABOUT ARRAYS

Essential: have items in predictable, and computable, locations in memory

=> "where is the i'th element"



Use predictable location to get from hash value to some specific index

=> Not all indices are used
=> Positions correspond to "array slots"

Use predictable location for get(i)

=> Items in consecutive locations in memory

USUALLY THINK ABOUT SHIFTING ELEMENTS TO MAINTAIN THIS

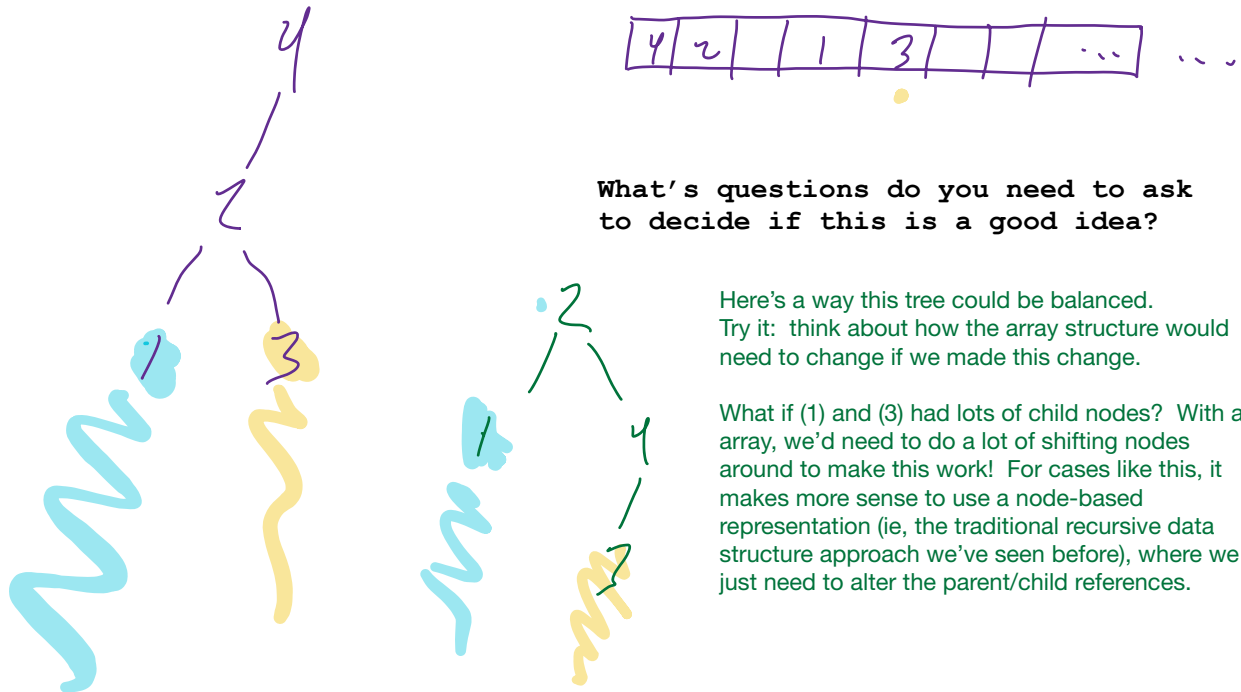
Use predictable location to navigate between parent/child nodes

=> Positions correspond to where we are in tree

Important to think about

- Ways different data structures can be used
- How the underlying data structure (arrays, in this case) can matter for different applications

Do all binary trees belong in arrays? What about BSTs?



Overall: think about how are going to use the data structure
=> What operations do you need to perform (at a high level)
=> For BSTs: need to keep the BST ordered and balanced

=> How does that translate into operations on the data structure
=> If we used an array, we would need to shift a lot of elements around to keep the BST balanced!

```

import math

"""implementation of a max heap"""
class Heap:
    def __init__(self):
        self.data = []
        self.size = 0

    def __str__(self):
        """string representation is the underlying list"""
        return str(self.data)

    def parent_index(self, of_index):
        """compute parent index of given index. Assumes of_index > 0"""
        return math.floor((of_index - 1) / 2)

    def swap(self, index1, index2):
        """swaps values in index1 and index2 within self.data"""
        tmp = self.data[index1]
        self.data[index1] = self.data[index2]
        self.data[index2] = tmp

    def insert(self, new_elt):
        """insert element into the heap"""
        self.data.append(new_elt)
        self.sift_up(self.size)
        self.size += 1

    def sift_up(self, from_index):
        """swap element in from_index up heap until it is in the right place"""
        if from_index > 0:
            parent = self.parent_index(from_index)
            if self.data[from_index] > self.data[parent]:
                self.swap(parent, from_index)
                self.sift_up(parent)

    def sift_up_while(self, from_index):
        """a while-loop based version of sift_up"""
        if from_index > 0:
            curr_index = from_index
            parent = self.parent_index(curr_index)
            while (curr_index > 0) and \
                (self.data[curr_index] > self.data[parent]):
                self.swap(parent, curr_index)
                curr_index = parent
            parent = self.parent_index(curr_index)
            # Note: this version repeats last two lines, unlike the recursive one

```

Heaps Implementation with Arrays